Improving Bidirectional Heuristic Search by Bounds Propagation

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Abstract

Recent research on bidirectional search describes anomalies, or cases in which improved heuristics lead to more node expansions. Aiming to avoid such anomalies, this paper characterizes desirable properties for bidirectional search algorithms, and studies conditions for obtaining these properties. The characterization is based on a recently developed theory for bidirectional search, which has formulated conditions on pairs of nodes such that at least one node from every pair meeting these conditions must be expanded. Moreover, based on this must-expand-pairs theory, we introduce a method for enhancing heuristics by propagating lower bounds (lb-propagation) between frontiers. This lb-propagation can bestow the desirable properties on some existing algorithms (e.g., the MM family) while avoiding the above anomaly altogether. Empirical results show that lb-propagation reduces the number of node expansions in many cases.

1 Introduction

Bidirectional heuristic search (Bi-HS) algorithms interleave two separate searches: a search forward from start, and a search backward from goal. Recently, a new line of research into Bi-HS was spawned. Eckerle et al. (2017) defined three conditions on the node expansions required by Bi-HS algorithms to guarantee solution optimality. Following work reformulated these conditions as a must-expand graph (GMM). It was shown that the Minimum Vertex Cover (MVC) of GMM corresponds to the minimal number of expansions required to prove optimality (Chen et al. 2017). Finally, a number of algorithms were introduced. NBS (Chen et al. 2017) and DVCBS (Shperberg et al. 2019) are non-parametric GMM-based Bi-HS algorithms that aim to find a vertex cover of GMM quickly, but in different ways. Fractional MM (fMM(p)) (Shaham et al. 2017) is a parametric algorithm that generalizes the MM algorithm (Holte et al. 2017) by controlling the fraction p of the optimal path at which the forward and backward frontiers meet. Another parametric algorithm, GFBHS (Barley et al. 2018), iteratively increases the depth of the search by using a split function to determine how deep to search on each side at each iteration.

Holte et al. (2017) observed an anomaly where improving a heuristic caused the MM algorithm to expand more nodes.

Aiming to generalize this anomaly beyond MM, Barley et al. (2018) defined that an algorithm is well-behaved if using a better heuristic will never hurt its performance; otherwise, it is ill-behaved. In this paper we expand this line of work on Bi-HS in several ways.

First, we study and develop desirable properties for Bi-HS algorithms by re-formalizing the well-behaved property and providing a definition which is even more general than that of Barley et al. (2018). We also introduce the reasonable property which guarantees that an algorithm will never expand nodes if the lower-bound associated with them is greater than the current global lower bound (LB) on the optimal solution. We then introduce and prove sufficient conditions required to fulfill each property.

Second, building on the conditions of Eckerle et al. (2017), we introduce lb-propagation, a method for propagating the best lower-bound between the two search frontiers, thereby improving heuristics and the f-values in each frontier. lb-propagation can be used on top of any Bi-HS algorithm; it is already used implicitly in GMM-based algorithms such as NBS and DVCBS. We show that lb-propagation causes the MM family to become well-behaved and reasonable, thereby avoiding the anomaly, although some algorithms, such as BS*, cannot be fixed in this way.

Third, we perform a study on a number of algorithms, characterizing those that are inherently well-behaved and reasonable, as well as whether or not lb-propagation bestows these properties on the algorithms. Finally, we show experimentally that lb-propagation reduces the number of node expansions for non-GMM-based algorithms.

1.1 Definitions and Background

A shortest-path problem, P, is defined as a tuple (G = (V, E), start, goal) in which G is a graph and start, goal ∈ V. The aim of such problems is to find the least-cost path between start and goal. Let d(x, y) denote the shortest distance between x and y and let C* = d(start, goal). In some cases, the minimal edge-cost is known beforehand; this minimal cost is denoted by ϵ.

Most Bi-HS algorithms maintain two open lists: OpenF for the forward search and OpenB for the backward search. There are two types of heuristics in bidirectional search. Front-to-front heuristics (de Champeaux 1983; de Champeaux and Sint 1977) estimate the distance between any
two nodes in the search space, while front-to-end heuristics (Kaindl and Kainz 1997) estimate the distance from any node and the start or goal. Front-to-front heuristics may be more computationally expensive, and efficient data structures for front-to-front algorithms do not exist. This paper only considers front-to-end heuristics.

Given a direction $D$ (either forward or backward) We use $f_D$, $g_D$, and $h_D$ to indicate $f$-, $g$-, and $h$-values in direction $D$. In addition, $f_{\text{min}}D$ and $g_{\text{min}}D$ represent the minimal $f$- and $g$-values in that direction.

The forward heuristic $h_F$ is admissible iff $h_F(u) \leq d(u, g)$ for every state $u \in G$ and is consistent iff $h_F(u) \leq d(u, u') + h_F(u')$ for all $u, u' \in G$. The backward heuristic $h_B$ is defined analogously. A pair of forward and backward heuristic functions is bi-admissible if both heuristics are admissible. Likewise, such a pair is bi-consistent if both heuristics are consistent. A search algorithm is admissible if it is guaranteed to find an optimal solution whenever its heuristic is admissible. Finally, a heuristic $h_1$ is said to dominate another heuristic $h_2$ if and only if for every node $n \in G$, $h_1(n) \geq h_2(n)$ (Russell and Norvig 2016). We limit the discussion in this paper to admissible deterministic black-box expansion-based algorithms (called DXBB by Eckerle et al. (2017)) used with bi-admissible and bi-consistent heuristics.

1.2 Fractional MM

We use the MM family of algorithms as a case study, therefore briefly describe them next. MM is a Bi-HS algorithm that meets in the middle (Holte et al. 2017), i.e. it is guaranteed to never expand a node whose $g$-value exceeds $C^*/2$. Fractional MM ($\text{fMM}(p)$) is a generalization of MM that never expands a node in the forward direction whose $g$-value exceeds $C^*/p$, and never expands a node in the backward direction whose $g$-value exceeds $C^*/(1-p)$. For a given fraction $0 < p < 1$, $\text{fMM}(p)$ chooses a node for expansion according to the following priority functions:

$$
pr_F(u) = \max(g_F(u) + h_F(u), \frac{g_F(u)}{p} + \epsilon) \\
pr_B(v) = \max(g_B(v) + h_B(v), \frac{g_B(v)}{1-p} + \epsilon)
$$

A node with minimal priority in either direction is chosen for expansion.1 $\text{fMM}$ terminates when one of the following conditions is met:

- One of $\text{Open}_F$ or $\text{Open}_B$ is empty.
- There exists a node $v$ in both open lists with $C = g_F(v) + g_B(v)$ s.t. either:
  - $f_{\text{min}}F \geq C$;
  - $f_{\text{min}}B \geq C$;
  - $g_{\text{min}}F + g_{\text{min}}B + \epsilon \geq C$; or
  - $\min\{ \min_{u \in \text{Open}_F} pr_F(u), \min_{v \in \text{Open}_B} pr_B(v) \} \geq C$.

Note that MM is a special case of $\text{fMM}(p)$ with $p = 1/2$.

Shaham et al. (2017) showed that for every problem instance, there exists a fraction $p^*$ such that $\text{fMM}(p^*)$ is optimally efficient and will expand the minimal number of nodes required to guarantee the optimality of its solution. However, $p^*$ is not known a priori since it depends on the search-tree structure and the value of $C^*$.

2 The Well-Behavedness Property

If $h_1$ and $h_2$ are consistent heuristics and $h_1(s) \geq h_2(s)$ for all non-goal nodes (i.e., $h_1$ dominates $h_2$), then every node expanded by $A^*$ using $h_1$ will also be expanded by $A^*$ using $h_2$ up to tie-breaking in the last $f$-layer (Holte 2010). Holte et al. (2017) describe an anomaly that may occur in Bi-HS algorithms such that a similar property does not hold. An example is provided in which MM using a global zero-heuristic (denoted henceforth by $h_0$ and the MM variant using it by MM$_0$) expands a subset of nodes that are expanded by MM that uses a stronger heuristic. Barley et al. (2018) also refer to the above anomaly, calling algorithms well-behaved when switching to a stronger heuristic does not lead to the expansion of any additional nodes, and ill-behaved otherwise. Well-behavedness has not been formally defined in a general manner; Holte et al. (2017) did not formally define the anomaly and Barley et al. (2018) defined it using terms that are specific to the GBFHS algorithm. We introduce a general definition of the well-behavedness property below and show that the anomaly results from a combination of (1) different tie-breaking, and (2) not using the theoretical lower-bound conditions for guiding the expansion process.

Many heuristic search algorithms do not fully specify which single node to expand at any given point in the search. For example, $A^*$ may choose any node in $\text{Open}$ with a minimal $f$-value, and $\text{fMM}$ can choose any node in either open list with minimal priority. Instead, these algorithms specify a set of nodes from the open lists (denoted henceforth by allowable-set) from which the next node must be expanded. An additional tie-breaking scheme is used to select a single node from the allowable-set. Tie-breaking is often specific to a given implementation, and in most cases is not part of the published algorithm definition. For example, $A^*$ must expand nodes with the smallest $f$-value. There are many possible tie-breaking rules to decide how to break ties among nodes with the same $f$-value (e.g., smallest $h$ or smallest $g$, generation order etc.). However, all of these tie-breaking functions are low-level details of $A^*$ implementations.

We use $A_h(I, t)$ to denote the sequence of nodes expanded by running algorithm $A$ using heuristic $h$ on problem instance $I$ with a tie-breaking function $t$, and by $S(A_h(I, t))$ the (unordered) set of nodes induced by the expansion performed by $A_h(I, t)$.

Definition 1. Let $h_1, h_2$ be bi-admissible bi-consistent heuristics, such that $h_1$ dominates $h_2$. Algorithm $A$ is said to be well-behaved if for every tie-breaking policy $t$ and problem instance $I$, there exists a tie-breaking policy $t'$ such that $S(A_{h_1}(I, t')) \subseteq S(A_{h_2}(I, t))$.

This is a general definition that can be used with any Bi-HS algorithm. To date, only $A^*$ and GBFHS have been proven to be well-behaved, while MM has been shown to be
2.1 Example of the Anomaly for $\mathbb{MM}$

In order to explore algorithms that are ill-behaved, we borrow an example from Holte et al. (2017), depicted in Figure 1. In this example $\epsilon = 1$ and the values inside nodes are $h$-values in the direction indicated by the arrow. We henceforth denote this bi-consistent heuristic by $h_{f_{fg}}$. $\mathbb{MM}_0$ expands nodes by their $g$-value. Thus, $\mathbb{MM}_0$ starts by expanding start and goal (priority of 0), after which nodes $S_1, G_1, A$, and $C$ have a priority of 1. There exists a tie-breaking policy $t$ in which $\mathbb{MM}_0$ expands $A, C$, and $S_1$ and then terminates, since it finds a solution of cost 4 and $g_{\text{min}} + g_{\text{min}} + \epsilon = 4$. In contrast, $\mathbb{MM}$ must expand both $S_1$ and $G_1$ after expanding start and goal before expanding $A$ and $C$, since $S_1$ and $G_1$ have a priority of $g + h = 2g + \epsilon = 3$, while $A$ and $C$ have a priority of 4 ($g + h = 4$). Consequently, there exists a tie-breaking policy $t$ for $\mathbb{MM}_0$ such that for every a tie-breaking policy $t'$ the set of nodes expanded by running $\mathbb{MM}$ on this instance using $t'$ is not a subset of the set of nodes expanded by $\mathbb{MM}_0$ using $t$. Thus, $\mathbb{MM}$ is ill-behaved.

To understand why $\mathbb{MM}$ is ill-behaved, consider the situation after $\mathbb{MM}$ expanded start, goal, and $S_1$. At this point, $\text{Open}_0$ contains $S_2$ ($g_F = 2, h_F = 1, pr_F = 5$), and $A$ ($g_F = 1, h_F = 3, pr_F = 4$); $\text{Open}_0$ contains $G_1$ ($g_B = 1, h_B = 2, pr_B = 3$), and $C$ ($g_B = 1, h_B = 3, pr_B = 4$). Thus, if the optimal solution goes through $G_1$, it must go through either $S_2$ or $A$. If the optimal path goes through $G_1$ and $S_2$, its cost would be at least $g_F(S_2) + g_B(G_1) + \epsilon = 4$. Similarly, if the optimal path goes through $G_1$ and $A$ then its cost would be at least $f_F(A) = 4$. Hence, every path that goes through $G_1$ must have a cost of at least 4. The priority of $G_1$ ($pr_F(G_1) = 3$) doesn’t reflect knowledge available in the search, which causes $\mathbb{MM}$ to be ill-behaved. This observation suggests that the sufficient conditions for node expansions (Eckerle et al. 2017) may be connected to whether an algorithm is well-behaved.

2.2 Guaranteeing Solution Optimality

Unidirectional search algorithms must expand all nodes $n$ with $f(n) < C^*$ in order to guarantee the optimality of solutions (Dechter and Pearl 1985).

Eckerle et al. (2017) generalized this to Bi-HS by examining pairs of nodes $(u, v)$ such that $u \in \text{Open}_F$ and $v \in \text{Open}_B$. Let $\epsilon$ be the minimal edge cost in $G^2$. If $u$ and $v$ meet the following conditions, then every algorithm must expand at least one of $u$ or $v$ in order to ensure that there is no path from $s$ to $g$ passing through $u$ and $v$ of cost $< C^*$.

1. $f_F(u) < C^*$
2. $f_B(v) < C^*$
3. $g(u) + g_B(v) + \epsilon < C^*$

Definition 2. For each pair of nodes $(u, v)$ let $lb(u, v) = \max \{f_F(u), f_B(v), g(u) + g_B(v) + \epsilon\}$

In Bi-HS, a pair of nodes $(u, v)$ is called a must-expand pair (MEP) if $lb(u, v) < C^*$. The MEP definition is equivalent to Eckerle’s conditions; for each MEP only one of $u$ or $v$ must be expanded. In the special case of unidirectional search, algorithms expand all the nodes with $f_F < C^*$, which is equivalent to expanding the forward node of every MEP. Bi-HS algorithms may expand nodes from either side, potentially covering all the MEPs with fewer expansions.

However, to address the ill-behavedness property we wish to bound the minimal solution cost that can pass through each node $u$ in our open lists. To do so, we use the bound $lb(u, v)$ and apply it to every node $v$ on the opposite frontier and take the minimum among these values. Formally, for every node $u$ in $\text{Open}_D$ let

$$lb(u) = \min_{v \in \text{open}_D} \{lb(u, v)\}$$

where $\text{D}$ denotes the opposite direction from $D$. Then, $lb(u)$ is a lower bound on the cost of any solution that passes through $u$. Finally, we define the global lower bound $LB$ to be the minimal $lb(u)$ among all nodes. This is identical to the minimal $lb(u, v)$ among all pairs. $LB$ was used in the high-level pseudocode (described below) of the NBS and DVCBS algorithms. Note that the search begins with $LB = lb$($start, goal$), after which $LB$ increases iteratively until $LB = C^*$. We can now use these definitions to show whether a family of algorithms is well-behaved.

2.3 Conditions for Being Well-Behaved

We first introduce three sufficient conditions for an admissible Bi-HS algorithm $A$ to be well-behaved:

Condition C1: Algorithm $A$ chooses a node $u$ for expansion only if $lb(u) = LB$.

Condition C2: Algorithm $A$ terminates as soon as a solution with cost $C \leq LB$ is found.

Condition C3: The allowable-set of algorithm $A$ contains every node $u$ with $lb(u) = LB$.

Theorem 1. An admissible Bi-HS algorithm $A$ that satisfies conditions C1, C2, and C3 is well-behaved.

Proof. Let $h_1, h_2$ be bi-admissible bi-consistent heuristics, s.t. $h_1$ dominates $h_2$, and let $l_2$ be an arbitrary tie-breaking

Strictly speaking the $\epsilon$ term was added by Shaham et al. (2018) as a generalization of the inequalities, since $\epsilon = 0$ is always a lower-bound to edge cost.
policy. We will show that there exists a tie-breaking policy \( t_1 \) s.t. \( S_1 = S(\text{A}_h(I, t_1)) \subseteq S(\text{A}_h(I, t_2)) = S_2 \).

To do this, we examine the execution of \( \text{A}_h(I, t_1 = t_2) \) and show how to modify \( t_1 \) to make \( S_1 \subseteq S_2 \). Let \( n \) be the first node expanded in the trace of the execution s.t. \( n \notin S_2 \) (if no such node exists then \( S_1 \subseteq S_2 \) and we are done). In addition, let \( D \) be the direction in which \( n \) was expanded.

We now have two cases:

**Case 1:** There exists a node \( n' \) in one of the frontiers s.t. \( lb(n') = lb(n) = LB \) and \( n' \in S_2 \), given \( C \), we can modify \( t_1 \) to choose \( n' \) instead of \( n \).

**Case 2:** All of the nodes \( n' \) in the frontiers with \( lb(n') = lb(n) = LB \) are not in \( S_2 \). We will show that this case is not possible. Note that since \( n \) was chosen for expansion, all other nodes \( m \) in the frontiers have \( lb(m) \geq lb(n) = LB \). Let \( v \) be a node in \( \text{Open} \) s.t. \( lb(n) = lb(v) \). Since \( lb(v) \leq lb(n) \) and \( lb(v) \geq lb(n) \), then \( lb(n) = lb(v) \), and therefore \( v \notin S_2 \). Since both \( n \) and \( v \) are in \( \text{Open} \) using \( h_1 \) when \( n \) is chosen for expansion, and because \( n \) is the first to not be in \( S_2 \), then all other nodes that were expanded by \( h_1 \) were expanded by \( h_2 \). Thus, at some point in \( \text{A}_h(I, t_2) \), both \( n \) and \( v \) are in \( \text{Open} \) with a \( g \)-value less than or equal to that found in \( \text{A}_h(I, t_1) \). We can therefore refer to \( lb(n, v) \) in \( \text{A}_h(I, t_2) \). Henceforth, \( lb(n, v) \) refers to the value \( lb(n, v) \) in \( \text{A}_h(I, t_1) \). Since \( h_1 \) dominates \( h_2 \), and since the \( g \)-values cannot be larger when using \( h_2 \), \( lb_2(n, v) \leq lb_1(n, v) \). We proceed by examining the possible values of \( lb_1(n, v) \), and show that any value leads to a contradiction. There are three possible values of \( lb_1(n, v) \) to consider:

1. If \( lb_1(n, v) < C^* \), then \( lb_2(n, v) < C^* \) as well. Therefore, \( \{n, v\} \) is an MEP and any admissible algorithm must expand one of them. The fact that \( n, v \notin S_2 \) contradicts the admissibility of \( A \).
2. If \( lb_1(n, v) = C^* \), then since \( \text{A}_h(I, t_1) \) expanded \( n \), it did not yet find any solution of cost \( C^* \). Since \( LB \) is a lower-bound on the optimal solution, a solution with cost \( C^* \) must pass through some node \( m \) with \( lb(m) = lb(n) = C^* \) that was not yet expanded. Under the assumption of case 2, there are no nodes \( m \) with \( lb(m) = lb(n) = LB \) in one of the frontiers that is also in \( S_2 \). Therefore, \( \text{A}_h(I, t_2) \) will never find a solution of cost \( C^* \). This is a contradiction to the assumption that \( A \) is admissible.
3. If \( lb_1(n, v) > C^* \), then since \( \text{A}_h(I, t_1) \) expanded \( n \) and \( A \) satisfies \( C_2 \), it did not find any solution of cost \( C^* \). Additionally, when \( n \) was chosen for expansion, the \( lb \) between every pair of nodes in the open lists is greater than \( C^* \). Thus, a solution of cost \( C^* \) does not exist by contradiction to the definition of \( C^* \).

These conditions are sufficient, but not necessary. In the next section we explore another desirable property.

### 3 The Reasonableness Property

While being well-behaved is an interesting property, some well-behaved algorithms do not behave sensibly. For example, an algorithm that completely ignores heuristic values and expands nodes according to their \( g \)-value is clearly well-behaved because a stronger heuristic will not change the behavior of the algorithm. However, such an algorithm might expand nodes \( n \) with \( f(n) > C^* \) whose \( g(n) \leq C^* \). Gilon, Felner, and Stern (2016) denoted algorithms as reasonable if they have a best-first structure (i.e. an open list and an expansion rule), and they prune any node \( n \) with \( f(n) > C \), where \( C \) an upper bound on the cost. We generalize this notion as follows:

**Definition 3.** A Bi-HS algorithm is reasonable if for every tie-breaking policy it does not expand a node \( v \) if either \( lb(v) > C^* \), or if \( lb(v) = C^* \) and a solution of cost \( C^* \) has already been found.

Note that since \( f(n) \leq lb(n) \) and \( C^* \leq C \), the redefinition of the reasonable property is tighter than the original definition of Gilon, Felner, and Stern (2016).

**Theorem 2.** Any admissible Algorithm \( A \) that satisfies \( C_1 \) and \( C_2 \) is reasonable.

**Proof.** Let \( A \) be an algorithm that always expands a node \( u \) with \( lb(u) = LB \) and terminates as soon as a solution with a cost \( c \leq LB \) is found. Assume by contradiction that \( lb(u) > C^* \). Since \( lb(u) \) is minimal \( lb(u) = LB \) then every solution that passes through every node in the open lists has a cost \( > C^* \). Since \( C_2 \) dictates that \( A \) terminates when a solution with a cost \( c = LB \) is found, no solution with cost \( C^* \) could have been found. Therefore, there is no possible solution with a cost of \( C^* \), by contradiction to the definition of \( C^* \).

To summarize both theorems, an algorithm that satisfies conditions \( C_1 \) and \( C_2 \) is reasonable, and one that also satisfies \( C_3 \) is well-behaved. In both cases, the conditions are sufficient but not necessary.

### 4 Improving Heuristics by \( lb \)-propagation

We next introduce several methods that improve the heuristic value of a node by utilizing information gathered during the search in both frontiers. The strongest method which propagates \( lb \)-values causes some ill-behaved algorithms to become well-behaved (e.g., the \( \mathbb{MM} \) family). In addition, algorithms that satisfy conditions \( C_1 \) and \( C_2 \) with respect to \( f \) instead of \( lb \) which use this method become reasonable. Note that this improvement is achieved by modifying only the heuristic, without any other changes to the algorithms.

#### 4.1 Propagating \( g \)- and \( f \)-values

A simple observation on the nature of bidirectional search yields that the minimum \( g \)-value with the addition of \( \epsilon \) is an admissible heuristic for any node in \( \text{Open}_D \). Furthermore, we can propagate the minimal \( f \)-value from the opposite frontier because it is a lower bound on any possible solution. Formally, let \( g_{\text{min}_D} = \min_{v \in \text{open}_D} \{ g(v) \} \) and let \( f_{\text{min}_D} = \min_{v \in \text{open}_D} \{ f(v) \} \). We can improve the heuristic of node \( n \) in direction \( D \) to be:

\[
    h'(D)(n) = \max\{ h(D)(n), g_{\text{min}_D} + \epsilon, f_{\text{min}_D} - g(D)(n) \}
\]

\( h' \) clearly dominates \( h \) and is easy to implement. One only needs to keep track of \( g_{\text{min}_D} \) and \( f_{\text{min}_D} \) for both directions. Nevertheless, \( h' \) does not solve the anomaly; \( \mathbb{MM} \) using
Table 1: Algorithm properties summary. R columns denote reasonableness, WB columns denote well-behavedness.

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Without lb-p</th>
<th>With lb-p</th>
</tr>
</thead>
<tbody>
<tr>
<td>BHPA</td>
<td>×</td>
<td>✓</td>
</tr>
<tr>
<td>BS*</td>
<td>×</td>
<td>×</td>
</tr>
<tr>
<td>IMM</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>GBFSH</td>
<td>✓</td>
<td>✓</td>
</tr>
<tr>
<td>NBS, DVCBS</td>
<td>✓</td>
<td>✓</td>
</tr>
</tbody>
</table>

Table 1: Algorithm properties summary. R columns denote reasonableness, WB columns denote well-behavedness.

\[ h' \text{ on Figure 1 behaves identically to } \text{IMM using the original heuristic, as described in Section 2.1} \]

### 4.2 lb-propagation heuristic

The next heuristic exploits knowledge from lb-values. Let \( h_{lb}(n) = lb(n) - g_D(n) \) denote the new heuristic function for nodes in direction \( D \). Consider the following key observations: (1) \( h_{lb} \) is a dynamic heuristic that takes into account information generated by the search in the opposite direction. Therefore, its value for a node may change as the search proceeds. (2) Since \( lb(n) \geq f_D(n), h_{lb}(n) \geq h_D(n) \) for every node in both directions. (3) \( h_{lb} \) maintains the bi-consistency and bi-admissibility properties of \( h \).

The heuristic \( h_{lb} \) dominates \( h' \) because \( h' \) looks at the global values of \( g_{min} \) and \( f_{min} \), while \( h_{lb} \) considers each pair of nodes in isolation. Despite the fact that \( h_{lb} \) dominates \( h' \), using lb-propagation depends on the ability to efficiently compute the \( lb \) of nodes in OPEN. This task is certainly more difficult than applying the other propagation, which simply requires maintaining the minimal \( f \)-and \( g \)-values in each direction. In some algorithms the lb-propagation can be applied to a limited subset of OPEN, possibly enabling an efficient implementation (similar to NBS). In other cases, the \( lb \) of every node is required. This leads to a potentially less a efficient implementation, using \( g \)-h buckets (Burns et al. 2012); this solution would work if the number of possible \( g \)-values (and therefore \( h \)-values) is relatively small, which is the case in many common domains.

An important property of \( h_{lb} \) is that it changes the \( f \)-values of nodes to be their \( lb \)-value, and therefore makes some existing algorithms well-behaved and reasonable as we show in the next section.

### 5 Classification of Existing Algorithms

As mentioned, lb-propagation makes the \( f \)-values of nodes identical to their \( lb \)-values. Therefore, any algorithm which chooses to expand nodes based on \( f \)-values and applies lb-propagation will now satisfy condition C1. However, in order to be provably reasonable it should also satisfy condition C2, and to be well-behaved, condition C3 is also needed. In this section, we review several Bi-HS algorithms and analyze how lb-propagation affects them. For any algorithm \( A \) we henceforth denote by \( A_{lb} \) a version of \( A \) that applies lb-propagation. Table 1 summarizes the results of this section, for algorithms with and without lb-propagation (lb-p).

#### 5.1 BHPA

We begin with BHPA (Pohl 1971), a simple algorithm that first selects a direction and chooses to expand a node with minimal \( f \)-value in that direction. BHPA terminates when the minimal \( f \)-value is greater than or equal to \( C^* \).

**Lemma 3. BHPA is well-behaved.**

**Proof.** Let \( I \) be a problem instance, \( h_1 \) and \( h_2 \) be heuristics that are bi-admissible and bi-consistent on \( I \) s.t. \( h_1 \) dominates \( h_2 \), and let \( t_2 \) be a tie-breaking policy. Let \( S_2 \) denote \( S(BHPA_{h_2}(I, t_2)) \) and let \( S_1 \) denote \( S(BHPA_{h_1}(I, t_1 = t_2)) \). Let \( u \) be the first node expanded in the trace of \( BHPA_{h_1}(I, t_1) \) s.t. \( u \notin S_2 \). If \( u \) does not exist, we are done. Otherwise, we want to fix \( t_1 \). If there exists a node \( n \in S_2 \) that has a minimal \( f \)-value in either direction of \( BHPA_{h_1}(I, t_1) \) when \( u \) was selected for expansion, we can alter \( t_1 \) to select \( n \) instead of \( u \). Otherwise, there exists a node \( u' \) in the opposite direction of \( u \) with minimal \( f \)-value, and we could modify \( t_1 \) to select \( u' \) instead of \( u \) for expansion. We know that for every node \( v, f_D^{h_1}(v) \geq f_D^{h_2}(v) \), thus \( f_D^{h_1}(u) \geq f_D^{h_2}(u) \) and \( f_D^{h_1}(u') \geq f_D^{h_2}(u') \). Therefore, if \( f_D^{h_1}(u) < C^* \) and \( f_D^{h_2}(u') < C^* \), we know that \( f_D^{h_2}(u') < C^* \) and \( f_D^{h_2}(u') < C^* \), hence \( S_2 \) must contain either \( u \) or \( u' \), so that \( BHPA_{h_2}(I, t_2) \) could terminate by contradiction to the fact that \( u, u' \notin S_2 \). Otherwise, \( f_D^{h_1}(u) = C^* \) or \( f_D^{h_2}(u') = C^* \). In this case, since \( BHPA_{h_1}(I, t_1) \) is admissible and must find an optimal solution, there must be some other node \( v \in S_2 \) in the open lists when \( u \) was chosen for expansion s.t. \( f_D^{h_1}(v) = C^* \) by contradiction to the case assumption. \( \square \)

**Lemma 4. BHPA is unreasonable.**

**Proof.** Consider the problem instance \( I \) in Figure 1 assuming that \( h_F(S_3) = h_B(G_3) = 0 \). Since for all \( i \in \{1, 2, 3\}, f_F(S_i) = f_B(G_i) = 3 \), while \( f_F(A) = f_B(C) = 4 \), running \( BHPA_{h_{lb}} \) on \( I \) with any tie-breaking policy must expand either \( \{S_1, S_2, S_3\}, \{G_1, G_2, G_3\} \), or both, before being able to expand \( A \) or \( C \). Furthermore, there exists a tie-breaking in which \( BHPA_{h_{lb}} \) expands \( start \) and \( goal \) followed by \( \{S_1, S_2, S_3\} \). Since \( lb(S_3) = g_F(S_3) + g_B(C) + \epsilon = 5 > C^* = 4 \), BHPA is unreasonable. \( \square \)

**Lemma 5. BHPA_{lb} is reasonable and well-behaved.**

**Proof.** Since after the propagation the \( f \)-value of a node equals its \( lb \), BHPA_{lb} always expands nodes with minimal \( lb \) \( (C1) \). In addition, BHPA_{lb} terminates as soon as a solution with \( \epsilon = 5 > C^* = 4 \). BHPA_{lb} is unreasonable. \( \square \)

#### 5.2 BS* (Kwa 1989)

BS* expands a node with a minimal \( f \)-value from the smallest open-list (Pohl’s cardinality criterion (Pohl 1971)) and terminates when the minimal \( f \)-value is greater than or equal to \( C^* \). In addition, BS* trims nodes from the open lists if their \( f \)-value is greater than or equal to costs of potential solutions that were already found.
Lemma 6. \( BS^* \) and \( BS^*_{lb} \) are ill-behaved.

Proof. Consider problem instance \( I \) in Figure 2. Using the heuristic values written inside the nodes, \( BS^* \) starts by expanding \( \text{start} \) and \( \text{goal} \) in an unspecified order. Nodes \( G_1 \) and \( S_2 \) get \( f \)-values of 2, and nodes \( S_1 \) an \( f \)-value of 5. Since \( BS^* \) has already found a solution of cost 4, nodes \( S_1 \) are trimmed from \( Open_F \). At this point \( Open_F \) contains only \( S_2 \), while \( Open_B \) contains two nodes (\( G_1 \)). Thus, \( BS^* \) is forced to choose \( S_2 \) for expansion before terminating.

Next, consider \( BS^* \) using \( h_B \). The beginning of the search is similar: \( \text{start} \) and \( \text{goal} \) in an unspecified order. Then all nodes of \( Open_F \) get an \( f \)-value of 1. Since no node is trimmed, \( Open_F \) contains 3 nodes, while the \( Open_B \) contains 2 nodes. Thus, \( BS^* \) expands the \( G_1 \) nodes before terminating, without expanding \( S_2 \).

While applying the propagation changes the \( f \)-value of nodes, the behaviour of \( BS^*_{lb} \) is identical to \( BS^* \) on this example. Thus, both \( BS^* \) and \( BS^*_{lb} \) are ill-behaved. We note that C3 is violated here because the allowable-set of \( BS^*_{lb} \) is forced to contain only one open list.

Lemma 7. \( BS^* \) is unreasonable.

Proof. The proof is similar to that of Lemma 4. Consider the problem instance \( I \) in Figure 1 assuming that \( h_F(S_3) = h_B(G_3) = 0 \). Similar to the proof of Lemma 4, \( BS^* \) will have to expand \( \text{start}, \text{goal}, \{S_1, S_2, S_3\} \) and \( \{G_1, G_2, G_3\} \) before expanding \( A \) or \( C \). Since \( lb(S_3) = g_F(S_3) + g_B(C) + \epsilon = 5 > C^* = 4 \), \( BS^* \) is unreasonable.

Lemma 8. \( BS^*_{lb} \) is reasonable.

Proof. Since after the propagation the \( f \)-value of a node equals its \( lb \), \( BS^*_{lb} \) always expands nodes with minimal \( lb \) (C1). Finally, \( BS^*_{lb} \) terminates as soon as a solution with a cost \( c \leq f_{min_D} = LB \) is found (C2). Therefore, both conditions are satisfied and \( BS^*_{lb} \) is reasonable.

5.3 \( fMM \)

We have already shown that \( fMM \) is ill-behaved in Section 2.1. We now show that \( fMM \) is also unreasonable.

Lemma 9. \( fMM \) is unreasonable.

Proof. Consider \( fMM(1/4) \) applied to the problem instance in Figure 3. After expanding \( \text{start} \) and \( \text{goal} \), a solution of cost 11 is discovered, and \( LB = lb(S_2, G_2) = g_F(S_2) + g_B(G_2) + \epsilon = 12 \). Since \( pr(G_2) = \max(f_B(G_2) + \frac{1}{2}g_B(G_2)) = 10 \), \( G_2 \) will be expanded before termination, even though 12 = \( LB > C^* = 11 \).

Figure 2: \( BS^*_{lb} \) is not well-behaved

Figure 3: \( fMM \) is not reasonable

Lemma 10. \( fMM_{lb} \) always expands a node \( u \) with \( lb(u) = LB \), hence it is reasonable.

Proof. Assume by contradiction that \( fMM_{lb} \) chose a node \( u \) in direction \( D \) for expansion s.t. \( lb(u') \neq LB \). Therefore, there exists a pair of nodes \((u', v')\) s.t. \( u' \) is in the \( Open_D \), \( v' \) is in \( Open_T \) and \( lb(u') = lb(v') = lb(u', v') = LB < lb(u) \). Using the \( lb \)-propagation, we know that \( f_D(u') = lb(u') = lb(u', v') \), and \( f_T(v') = lb(v') = lb(u', v') \). Therefore, \( f_D(u) = lb(u) > f_D(u') \). Likewise, \( f_D(u) = lb(u) > f_T(v') \).

\[ f_D(u) > f_T(v') = f_D(u') = lb(u') = lb(u', v') \]
\[ \geq g_D(u') + g_T(v') + \epsilon \]

Since \( u \) was chosen for expansion, we know that \( pr_D(u) \leq pr_D(u') \) and \( pr_D(u) \leq pr_T(v') \). Thus,
\[ \max(f_D(u), \frac{g_D(u')}{p} + \epsilon) \leq \max(f_D(u'), \frac{g_D(u')}{p} + \epsilon) \]
and
\[ \max(f_D(u), \frac{g_D(u)}{p} + \epsilon) \leq \max(f_T(v'), \frac{g_T(v')}{1-p} + \epsilon) \]

Since \( f_D(u) = lb(u) > lb(u') = f_D(u') = lb(v') = f_T(v') \),
\[ f_D(u) \leq \frac{g_D(u')}{p} + \epsilon \] (1)
\[ f_D(u) \leq \frac{g_D(u')}{1-p} + \epsilon \] (2)

By summing inequality (1) multiplied by \( p \) with inequality (2) multiplied by \( 1 - p \), we get:
\[ f_D(u) = p f_D(u) + (1-p) f_D(u) \leq g_D(u') + g_T(v') + \epsilon \]
In contradiction to: \( f_D(u) > g_D(u') + g_T(v') + \epsilon \) above.

Lemma 11. \( fMM_{lb} \) (with \( lb \)-propagation) is well-behaved.

Proof. Here we cannot use Theorem 1 directly since the allowable-set of \( fMM_{lb} \) does not include every node \( u \) with \( lb(u) = LB \), as some of these nodes might have \( g \)-values that raise their priority. Nonetheless, we show that \( fMM_{lb} \) using \( lb \)-propagation is in fact well-behaved, using a slight modification to Theorem 1. In Lemma 10 we showed that \( fMM_{lb} \) only expands nodes with minimal \( lb \). In addition, \( fMM_{lb} \) terminates when the lowest \( f \)-value is \( \geq C^* \). Since the \( f \)-value of nodes after applying \( lb \)-propagation is equal their \( lb \)-value, \( fMM_{lb} \) stops when \( LB \geq C^* \). Thus, C1 and C3 are satisfied. However, C2 is violated by the priority mechanism of \( fMM \), since nodes with the same \( lb \)-value
might have different priorities due to their $g$-values and direction. For example, if there is only one node in $Open_F$ with $f_B = 3, g_F = 1$, and only one node in $Open_B$ with $f_B = 3, g_B = 2$, $f_{MM}$ will give lower priority to the node in $Open_F$, based on his $g$-value (priority of 3 versus a priority of 5). Nonetheless, we will show that $f_{MM}$ is still well-behaved. The problem arises since the first case of the proof of Theorem 1 reduces to nodes $n'$ with the $lb(n') = lb(n)$ and $pr(n') = pr(n)$. Therefore, nodes $n'$ with $lb(n') = lb(n)$ and $pr(n') \neq pr(n)$ are part of the second case of the proof, which does not cover them. In that case, we considered some $v$ s.t. $lb(n) = lb(v) = lb(n, v)$. Following that, it was clear that $v \notin S_2$. Nonetheless, in our case, $v$ could have been in $S_2$ if it had a different priority than $n$ when $f_{MM}$ was running using $h_1$. Since $f_{MM}$ expands nodes with minimal priority, $pr(v) > pr(n)$. The priority of $v$ could have been determined by one of the following options:

**Case 1:** $f_D(v) = lb(v) = pr(v)$. However, $lb(u) \leq pr(u)$ and $lb(v) = lb(u)$. Therefore, $pr(v) \leq pr(u)$, by contradiction to the assumption that $pr(v) > pr(u)$.

**Case 2:** $g_D(v) = pr(v)$. Since $f_D(v) \leq f_D(v)$ and $f_D(v) \leq g_D(v) = pr(v)$, we know that the $pr(v)$ using $h_2$ is less or equal than $pr(v)$ using $h_1$. In addition, the priority of node $n$ (and any of its ancestors) when running using $h_1$ must be strictly less than $pr(v)$, by contradiction to the fact that $v$ was already chosen for expansion.

Therefore, we can conclude that $v$ is still not in $S_2$ and the proof of Theorem 1 is generalized to $f_{MM}$ as well.

### 5.4 NBS and DVCBS

NBS (Chen et al. 2017) and DVCBS (Shperberg et al. 2019) are two prominent Bi-HS algorithms that choose nodes for expansion with minimal $lb$. At any point in the search, NBS chooses a pair of nodes with minimal $lb$ and expands them both. DVCBS expands nodes from a subset of those with minimal $lb$, determined by maintaining a dynamic version of the $G_{MX}$ (denoted by $DG_{MX}$) and finding its MVC. Both algorithms terminate as soon as a solution with a cost $\epsilon \leq LB$ is found. Since the expansion policy and termination condition of both algorithm already consider $LB$, their properties remain unaffected by $lb$-propagation.

Clearly, NBS and DVCBS satisfy conditions C1, and C2 and are therefore reasonable (up to a single additional expansion). However as previously mentioned, the allowable set of DVCBS includes only nodes that make up the MVC of $DG_{MX}$, violating C3. In addition, once NBS has chosen a pair $(u, v)$ for expansion, it is committed to expanding both nodes. Therefore, after expanding $u$, any node $u'$ in the same direction of $u$ s.t. $lb(u') = lb(u) = lb(v) = LB$ is not in the allowable set of NBS until after expanding $v$, in violation of condition C3.

**Lemma 12.** NBS is ill-behaved.

**Proof.** Consider the problem instance of Figure 4.

Figure 4: NBS is not well-behaved.

wards, $X$ and $A$ are added to $Open_F$, and $C$ is added to $Open_B$. Since $g_F(x) = f_F(x) = 12, lb(X, C) = 12$, while $lb(A, C) = 9$. Therefore, the pair $(A, C)$ is chosen for expansion, after which a path of length 20 has been discovered. Since $lb(X, B) \geq 20$ and $lb(B, B) \geq 20$, NBS terminates after expanding $start, A, C$, and $goal$.

NBS using the heuristic in the nodes starts by expanding $start$ and $goal$. Next, $X$ and $A$ are added to $Open_F$, and $C$ is added to $Open_B$. Since $lb(A, C) = 20$ (due to $f_F(A)$) and $lb(X, C) = 19$ (due to $f_B(C)$). Therefore, NBS will expand the pair $(X, C)$, despite the fact that $X$ was not expanded using a weaker heuristic.

**Lemma 13.** DVCBS is ill-behaved.

**Proof.** Consider the problem instance of Figure 2. DVCBS using the heuristic in the nodes starts by expanding either $start$ or $goal$ since both of them are MVCs of $DG_{MX}$. If $goal$ was chosen, $start$ must be expanded, since it becomes the only MVC, followed by $S_2$ for a similar reason, after which DVCBS terminates since $LB \geq 5 > 4$ (the path that was discovered from $start$ to $goal$). Likewise, if $start$ was chosen for expansion, DVCBS will expand either $S_1$ and terminate, or $goal$ followed by $S_1$. In both cases the $G_1$ nodes are never expanded. However, DVCBS using $h_0$ must expand $start$ and $goal$ in an unspecified order, followed by $G_1$.

### 5.5 GBFHS

GBFHS is an algorithm that iteratively increases the depth of the search ($f_{Lim}$). At each depth, a pre-defined split function (parameter of the algorithm) is used that determines how deep to search on each side at each iteration by splitting $f_{Lim}$ to $g_{Lim_F}$ and $g_{Lim_B}$ s.t. $f_{Lim} = g_{Lim_F} + g_{Lim_B} + \epsilon - 1$. At every iteration, GBFHS considers nodes for expansion in direction $D$ with $f \leq f_{Lim}$ and $g < g_{Lim_D}$. GBFHS terminates when as soon as a solution with a cost equals to $f_{Lim}$ is found.

GBFHS was proved to be well-behaved. We show that GBFHS is reasonable by showing that it expands only nodes with minimal $lb$, even without applying $lb$-propagation.

**Lemma 14.** GBFHS is reasonable.

**Proof.** Since GBFHS considers nodes for expansion in direction $D$ with $f \leq f_{Lim}$ and $g < g_{Lim_D}$, a node $u$ that is chosen for expansion will have $lb(u) \leq \max\{g_{Lim_F} + g_{Lim_B} + \epsilon - 1, f_{Lim}\} = f_{lim}^5$, and since the $f_{lim}$ is increased only after there are no nodes left for expansion,

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5. Even though well-behavedness was not defined in a general manner when GBFHS was created, the proof of Barley et al. (2018) is still applicable to the new definition with slight modifications.

Barley et al. (2018) implicitly assume that $\epsilon$ is an integer $\geq 1$. 

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\footnote{This example is due to Robert Holte and Sandra Zilles}
\( lb(u) = f_{lim} = LB \). Ergo, GBFHS expands nodes with minimal \( lb \). In addition, GBFHS terminates as soon as a solution of cost \( f_{lim} = LB \) is found. Thus, C1 and C2 are satisfied and GBFHS is reasonable.

### 6 Experimental results

We ran experiments on three domains: (1) 50 10-Pancake Puzzle instances with the GAP heuristic (Helmer 2010). To get a range of heuristic strengths, we also used the GAP-\( n \) heuristics (for \( n = 1 \ldots 9 \)) where the \( n \) smallest pancakes are deleted from the heuristic computation; (2) 50 instances of the 10-disk 4-peg Towers of Hanoi (TOH4) problem with (8+2) and (6+4) additive PDBs (Felner, Korf, and Hanan 2004). (3) Grid-based pathfinding: 65 maps from Dragon Age Origins (DAO) (Sturtevant 2012), each with different start and goal points (a total of 1,680 instances).

Figure 5 shows the average number of nodes expanded by \( \text{MM} \) and by \( \text{MM}_{lb} \) in the 10-pancake domain across all GAP heuristics. Clearly, adding the \( lb \)-propagation significantly reduces the number of nodes expanded. Using \( h_{lb} \) seems to reduce the number of node expansions for each of the GAP heuristics up until GAP-7, in which the heuristic effectively becomes \( h_0 \). In addition, this figure clearly demonstrates the anomaly of \( \text{MM} \); the average number of nodes expanded by \( \text{MM} \) using heuristics GAP-2 through GAP-6 is greater than the number of nodes expanded by using heuristics GAP-7 through GAP-9 (notice the “hump-in-the-middle” (Barley et al. 2018)). By contrast, the hump-in-the-middle of \( \text{MM}_{lb} \) is much smaller, and in fact not visible when considering the average number of expansions. However, there were still some individual problem instances in which \( \text{MM}_{lb} \) expanded fewer nodes using a weaker heuristic. This is consistent with Theorem 1, since we are using a predetermined tie-breaking policy and not the best possible tie-breaking policy for every instance. Interestingly, \( \text{MM}_{lb} \) using \( \epsilon = 0 \) demonstrates no hump-in-the-middle, even when considering individual problem instances.

Similarly, Figure 6 shows the average number of nodes expanded in the 10-pancake domain across all GAP heuristics, with \( \epsilon = 1 \) by a variant of BHPA denoted by BHPA-Min. BHPA-Min selects the frontier that includes the node with the minimal \( f \)-value. Here too, the \( lb \)-propagation improves the search by reducing the number of nodes expanded. Even though BHPA-Min is well-behaved, and demonstrates no hump-in-the-middle in the average case, the \( lb \)-propagation still improves the algorithm by making it reasonable. This improvement is more evident with GAP-8 and GAP-9; Despite these GAP heuristics behaving like \( h_0 \) in the 10-pancake domain, \( lb \)-propagation incorporates \( g_{min_F} + g_{min_B} + \epsilon \) into the \( f \)-values of nodes in BHPA-Min, exposing an additional termination condition, and allowing the search process to halt sooner.

The average number of node expanded across domains, using \( \epsilon = 1 \), appear in Table 2. There is one row for each algorithm, and one column for each of the domains and their heuristics; \( h \) denotes the original heuristic, while \( h_{lb} \) denotes the heuristic enhanced by \( lb \)-propagation. The algorithms we have tested are \( BS^* \), \( fMM(p) \) using \( p \in \{1/4, 1/2, 3/4\} \), BHPA-Min and BHPA-Alt, another variant of BHPA that alternates between the frontiers between expansions. The results show that using the \( lb \)-propagation reduces the number of node expansions in most cases by up to a factor of 4. The \( lb \)-propagation particularly excels when the heuristics are weak. In these cases using \( h_{lb} \) always result in fewer node expansions; this is also the case for GAP-4 through GAP-9, which do not appear in the table. Another interesting observation is that the hump-in-the-middle is less pronounced in all tested algorithms. \( BS^* \) seems to be the least affected by the propagation among all algorithms. We posit that the reason for this is that \( BS^* \) is highly dependent on the search process; since \( BS \) selects a node for expansion from the direction with the smallest open list, minor increments in heuristic values might cause nodes to be trimmed away, possibly changing the size balance of the two frontiers. \( BS^* \) also assumes that the heuristic is consistent, which other algorithms do not. We have also experimented using \( \epsilon = 0 \); the results are similar to the reported \( \epsilon = 1 \) results.

Naturally, maintaining and using \( lb \) for each node incurs overheads. In the domains we used, \( g \)-bucketing requires negligible time and space. However, we did not focus on code optimization, and used naive data-structures. Thus, run times are not reported here. Improving efficiency with a bucketing scheme or adapting the data structures used for efficient \( lb \) computations by \( NBS \) is reserved for future work.
Table 2: Experimental results of average node expansions across domains

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>10-Pancake</th>
<th>TOH-10</th>
<th>Grid</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>8+2</td>
<td>64+4</td>
<td></td>
</tr>
<tr>
<td>BPHA-Alt</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>BPHA-Min</td>
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<tr>
<td>BS*</td>
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<tr>
<td>MM(1/4)</td>
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</tr>
<tr>
<td>MM(1/4a)</td>
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</tr>
<tr>
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<td>26</td>
<td>374</td>
</tr>
<tr>
<td>BPHA-Min</td>
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<td>25</td>
<td>374</td>
</tr>
<tr>
<td>BS*</td>
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<td>25</td>
<td>682</td>
</tr>
<tr>
<td>MM(1/4)</td>
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<td>115</td>
<td>6,002</td>
</tr>
<tr>
<td>MM(1/4a)</td>
<td>64</td>
<td>81</td>
<td>6,002</td>
</tr>
</tbody>
</table>

7 Discussion

We have examined the source of the anomaly exhibited by some Bi-HS algorithms, where using a better heuristic causes the algorithm to expand more nodes. Aiming to improve some algorithms in which the anomaly manifests, the properties of “well-behavedness” and “reasonableness” were defined, and sufficient conditions (C1, C2, C3) for these properties were established. These properties provide insights that lead to the lower-bound propagation scheme (lb-propagation) that can be added to many existing Bi-HS algorithms, in some cases bestowing upon them these desirable properties. Empirical results show that modified algorithms exhibit better behavior, alleviating or even eliminating the undesirable “hump-in-the-middle” effect seen when an algorithm is run with heuristics of varying quality.

The well-behavedness property as defined in this paper ensures that there exists a tie-breaking policy for which the anomaly would not occur. However, the desired tie-breaking policy is not specified. It is a non-trivial issue, left for future research, to define conditions that guarantee a stronger well-behavedness property of an algorithm, such that a dominating heuristic would never cause more nodes to be expanded than the weaker heuristic using the same tie-breaking policy. Another interesting research direction is to re-examine the three well-behavedness conditions with respect to heuristics that are strictly dominating, i.e., $h_1 > h_2$.

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